

Amendments to the Specification:

Please replace the paragraph on page 7, lines 2-3, with the following amended paragraph:

FIGs. 12 to 20 comprise pages of a flowchart showing how the video file ~~system~~ server of FIG. 1 is programmed to respond to a client's request for video service.

Please replace the paragraph on page 10, line 18, to page 11, line 9, with the following amended paragraph:

It is possible to configure the hardware of the video file server 20 for effective service of all kinds of movies due to a certain regularity in the frequency of access of movies once the movies have been ranked in terms of their frequency of access. Shown in FIG. 2, for example, is an idealized plot of the frequency of access, in plays per day, as a function of the ranking of each movie accessible from the video file server. The most popular movie is given a rank of one, the next most popular movie is given a rank of 2, and so on. When plotted on log-log paper, the plot falls on a straight line with negative slope in accordance with the "Zipf" distribution. (See, for example, the Jakob Nielsen, "Zipf Curves and Website Popularity," published at <http://www.useit.com/alertbox/zipf.html>). Although the popularity of each movie may rise and will fall with time, the plot will retain its general characteristic of a line having a certain slope and intercept with the "y" axis. Therefore, it is possible to configure and program the video file server in accordance with the general characteristic, in order to handle effectively each movie

ranking. Moreover, the video file server can be configured and programmed initially without regard to the rank of any particular movie.

Please replace the paragraph on page 14, lines 1-10, with the following amended paragraph:

For the most popular movies, there can be multiple data movers servicing just one movie, and in this case only a fraction of the cache RAM for servicing an entire movie need reside in any one of the movies. For example, as shown in FIG. 4, the movie data for movie #1 could be circulated among the first three data movers, and the movie data for movie #2 could be circulated among the third and fourth data movers. The data mover #1 could have a cache capacity to hold 40% of movie #1, the data mover #2 could have a cache capacity to hold 40% of movie #1, the data mover #3 could have a cache capacity to hold 25% of movie #1 and ~~40%~~ 40% of movie #2, and the data mover #4 could have a cache capacity to hold 65% of movie #2. Data mover #6 could have a cache capacity to hold at least 100% of movie #4.

Please replace the paragraph on page 17, lines 1-20, with the following amended paragraph:

FIGs. 12 to 20 comprise pages of a flowchart showing how the video file ~~system~~ server of FIG. 1 is programmed to respond to a client's request for video service. This request is received by the control station (22 in FIG. 1). In a first step 71, the control station (based on initial state in the service data base 32 in FIG. 1) selects a data mover DM_j to serve the client.

The control station sends a message over the high-speed serial bus (31 in FIG. 1) requesting DMj to respond to the client request. Then in step 72, the data mover DMj connects to the client. In step 73, the client negotiates a movie title with the data mover DMj, for example, as described above with reference to FIGS. 10 and 11, so that the data mover DMj obtains from the client a title of a movie that has been placed in disk storage. In step 74, the data mover DMj checks whether it has access to the movie. For example, DMj must be in the set of data movers assigned to service the rank of the movie. It is also possible that DMj is in the set of data movers assigned to service the rank of the movie but the movie should be serviced by another one of the data movers in the set, for example, because the portion of the movie to be initially accessed presently resides in the local cache of another one of the data movers in the set. If the titled movie is not accessible by the data mover DMj, then execution branches to step 75. In step 75, DMj informs the control station (CS) of the movie title. Then in step 76, the control station selects a data mover having access to the titled movie, and this selected data mover will be referenced as DMj in subsequent steps. Execution then continues to FIG. 13. Execution also continues to FIG. 13 from step 74 if the title is accessible by DMj.